

### **OVS Acceleration using Network Flow Processors**

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### Agenda

- Background: OVS on Network Flow Processors
  - Network device types => features required => acceleration concerns
- OVS Acceleration Options
  - OVS (or OpenFlow) Agent Only
  - Offloading OVS with Fallback
    - Examples:
      - Userspace <=> accelerator
      - Userspace <=> kernel <=> accelerator
      - Userspace <=> (kernel OR accelerator)
      - Userspace <=> (kernel AND accelerator)
- Observations
- Evolution
- Conclusions





**Network Flow Processor Powered Networking Device** 























### Any app on server

Middlebox apps:

Protocol

OpenFlow Controller and/or other control /

management software







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OpenFlow







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Middlebox apps:

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OpenFlow

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### OVS Acceleration Approaches: OVS Agent vs OVS Offload/Fallback

### "OVS = Agent, NFP = Switch"





- Fallback reasons: table lookup miss (table overflow vs. not full but no entry matched), unsupported classification (e.g. unsupported protocol), unsupported action, and/or deliberate hand-off to agent (e.g. packet in to controller)
- Most appropriate approach depends on features / perf. needed (related to type e.g. NIC in server, intelligent TOR, mbox / gateway)
- In part depends on performance required vs. achievable with attached x86 (or other CPU)

All variations support packet delivery from NFP directly to (host / guest) x (kernel / user mode), tunnel handling in NFP...

# "Offload to NFP - Fallback to x86"



### "OVS Agent Only" Features

### Dataplane feature set "limited" to what switch/NIC hardware (NFP) supports

- Could exceed what OVS dataplane offers (e.g. set individual header fields, not just entire L2/L3 header like OVS kernel datapath)
- Netronome's implementation actually fairly complete (with high performance):
  - OpenFlow 1.3+ with 250 tables, Ethernet / VLAN / MPLS / IPv4+v6 / TCP / UDP... matching + actions (set/push/pop/dec TTL), metering, QoS, logical port style tunnel termination/origination (VXLAN / [NV-]GRE / ...)
  - Extensions: fragmentation, ICMP stack, IP stack (route lookup / ARP processing) for tunnels, load balancing, improved QoS...
- Approach can support (host or guest) x (user or kernel mode) data delivery (incl. SR-IOV)
  - Apps (kernel netdev / user mode driver) modeled as virtual ports after app transited, packets re-injected into pipeline
- Apps (user mode driver) can trigger egress fastpath after app transited, packet sent directly to egress port
- Load balancing to application instances (static + stateless vs. dynamic + stateful)
- Each physical switch/NIC port represented by a VNIC instance in x86 (typically host kernel netdev)
- Exposes physical port statistics, link state -- Linux commands e.g. ethtool / ip / ifconfig just work
- In-band control: my MAC traffic diverted, broadcast / multicast copied (without needing OpenFlow table entries)
- Permits sniffing traffic (tcpdump) detects promiscuous mode on/off (could offload BPF too)
- Flow tracking table
- Manual population / modification via API (n M mods/s) or OpenFlow (mods/s are limited by controller / agent)
- Auto learning n M microflows/s optionally using API to update policy / obtain statistics



### Offloading with Fallback: Offload OVS Kernel vs. Userland



- Fallback to OVS x86 code to support features not in NFP (yet) or entries not present in NFP tables (e.g. DRAM-less NIC)
- OVS kernel offloaded to NFP, fallback to OVS kernel (and fallback onward to OVS userland as usual)
  - Feature set typically reflects (but can sometimes exceed) OVS kernel datapath capabilities
  - Could be regular OVS userland but also different userland (BigSwitch, Midokura, PLUMgrid...)
- OVS userland offloaded to NFP, fallback to OVS userland (via zero copy driver e.g. DPDK)
  - OVS kernel datapath not involved / required (no impact on features + performance, fewer issues with upstreaming to kernel.org)
- Either way a VNIC instance (e.g. kernel netdev) per physical port is useful statistics, link state, sniffing, in-band control (OpenFlow / SSH ...)

All variations support packet delivery from NFP directly to (host / guest) x (kernel / user mode), tunnel handling in NFP...

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### Fallback with Direct Host / Guest Delivery



- NFP sends fastpathed (non-fallback) packets directly to (host or guest) x (kernel or userspace) e.g. to SR-IOV VF VNIC3
- Fallback traffic sent to host (kernel / userspace): presented to OVS via e.g. VNIC OFP1' --- say OVS outputs to VNIC3'
- Reflector forwards from VNIC3' to VNIC3 (no need to re-process in pipeline as already processed by fallback code)

rspace) - e.g. to SR-IOV VF VNIC3 ' --- say OVS outputs to VNIC3' processed by fallback code)



### Offloading with Fallback: Combos

- Hybrid OVS userland / kernel offload + fallback
- Easy: multiple switch instances: some userland, others kernel each physical port attached to one of these
- Harder: within a switch instance, some traffic sent to kernel, other traffic to userland
  - Based on what?
    - Determination whether userspace processing will be needed if so, skip kernel?
      - Does not make sense if userspace fallback is minority of traffic / userspace is slower than kernel
    - Default to userspace, change to kernel when kernel processing is needed (for e.g. conntrack, IPsec, etc.)
      - Useful if userspace is faster than kernel
  - Is determination accurate?
    - Hopefully easy and accurate for entire microflow if based on traffic type (stateless)
    - Handing over or starting mid-flow can causes issues (entry state missing in tables, statistics wrong etc.)
  - Not urgent to implement on high functionality / capacity / throughput platform like NFP as most traffic is handled by it...
- Hybrid SDN / traditional networking, hybrid OpenFlow and non-OpenFlow features
- Traditional forwarding = L2/L3
  - Again, easier if ports attached to distinct traditional or switch instances (internal ports to link them)
  - Otherwise need one mechanism (SDN/traditional) to be primary, other secondary, or need third classifier as tie-breaker
- Non-OpenFlow features like tunnel handling (e.g. VXLAN / IPsec) / firewalling / NAT / other "black box" features
  - Invocation of features via built in behavior (e.g. my MAC / traffic type) vs. explicitly via an OpenFlow action / logical port etc.
  - Some need IP stack (ARP, route lookup, frag tracking, fragment / defragment) or other kernel functionality

### Examples as implemented on NFP

- Logical port style VXLAN / GRE termination handled via OVS kernel style mechanisms (lookup tunnel header + inner header together in tables)
- IP stack (route lookup to obtain egress port + source IP) and ARP processing for tunnel origination, leveraging Linux stack (entries are cached on NFP)
- IKE for IPsec, leveraging Linux usermode code
- Fragment tracking (without reassembling fragments) + peeking into tunnels (without terminating them) + flow tracking -- implemented "before" OVS / OpenFlow

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### **Observations**

- NFPs support all these acceleration approaches (implemented agent, usermode offload, kernel offload)
- Certain approaches may be more or less suited to more limited acceleration hardware
- Offload with fallback degrades gracefully (features e.g. classification / actions, also capacity)
- Fallback to kernel vs to usermode?
  - Where are the existing standard features (kernel e.g. conntrack, vs. usermode e.g. richer OpenFlow)?
  - · Which parts of the code do OVS community customize most often / readily?
  - Which is most performant (e.g. DPDK or other userspace driver avoiding some kernel overheads vs. cache in kernel)?
- Offloading OVS kernel does not necessarily yield the highest performance examples of issues:
- OVS kernel datapath functionality limited, e.g. broad brush actions: decrement TTL => replace L3 header
- Number of entries in OVS kernel tables could explode or experience churn (even with megaflow changes)
- Offloading OVS userspace simplifies supporting acceleration hardware with varying intelligence / capacity
- In kernel / at netlink (DPIF), entries have been "compiled" to low level => easier when seeing higher level intent in user mode
- Deciding where to perform processing (userspace / kernel / accelerator, OVS specific code / standard Linux) kernel code etc.) could be complex
- NFP's OVS kernel offload is fully featured, has large table capacity etc. => easy to offload all kernel table entries with small kernel patch
- For more limited devices, decision making may be more complex functionality split => table entry positioning best handled in userspace
- Can still decide in userspace to offload directly to accelerator, vs. via kernel to accelerator



## **Evolution of Dataplane Flexibility**



x86 Apps Calling APIs Compatible / Open Sourced vs. Vendor Extension



Flow API Load Balancing API Crypto API Forwarding API RDMA API ...

**Datapath Extensions** in NFP Native Code (e.g. C)











### **Flexible Datapath Abstraction**

OpenFlow 2.x, P4, PIF, eBPF...

Protocol agnostic flexible parsing

Arbitrary arrangement of matching tables

> Matching without tables

State storage / retrieval

Complex actions

Event handling



### Conclusions + Next Steps

- Performance examples (details depend on traffic patterns, number and type of flow entries / actions etc.)
- Using "20G" NFP, measured >20x speedup vs. unaccelerated OVS (IP + MPLS forwarding use case, multi table)
  - Difference can be larger for more complex actions / tunnels / traffic patterns e.g. VXLAN, IPsec, high flow setup rates
  - Observed flow tracker performing flow learning (5-tuple) at 12Mfps
- Expect further improvement (~3x-10x) for "200+G" NFP
- Capacities: millions or tens of millions of flow entries, 100,000s of tunnels (requirements vary per device type)
- => Order(s) of magnitude improvement achievable using acceleration
- Questions for OVS community
- Where to focus going forward implement features in kernel (for which OSes) vs. implement in userspace (cross OS / lightweight) vs. both
- Considerations: TTM, software only performance, ease of acceleration, leveraging existing code / developer skill sets, ease of maintenance
- Questions for Linux and other OS communities
- Leverage OVS vs implement different mechanisms e.g. tunnel termination/origination, QoS, eBPF / PIF...
- Questions for all
- How best to support acceleration hardware, without duplicating efforts
- (Excluded due to lack of time: specific API proposals for table manipulation APIs, acceleration "objects" and APIs: Linux/OVS/other...)

